

Practical Conversion from CPS to Direct Style

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CPS is great for compilers

- Evaluation order is made explicit.
- Control-flow is regularized.
- Useful for both high-level and low-level representations.
- Easily supports non-local control-flow; exceptions, call/cc, *etc.*

Bringing Continuations to LLVM

- Ongoing work to explore implementations of continuations.
- Native codegen is a pain; using LLVM is easier.
- Recent work: heap-allocated, first-class conts with LLVM
- How can a CPS-based compiler use LLVM with a stack?

CPS with a stack

```
fun fib (n, k) =  
  if n <= 2 then k 1  
  else  
    cont minus2 f_1 =  
      cont add f_2 =  
        k (f_1 + f_2)
```

minus2's closure
dies here

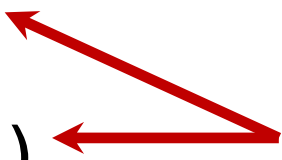
in

fib (n-2, add)

in

fib (n-1, minus2)

Downward funargs



Undoing CPS *in theory*

Key Observation*

Most continuations created by CPS are well-behaved.

* by Danvy, Kelsey, *etc.*

Undoing CPS *in practice*

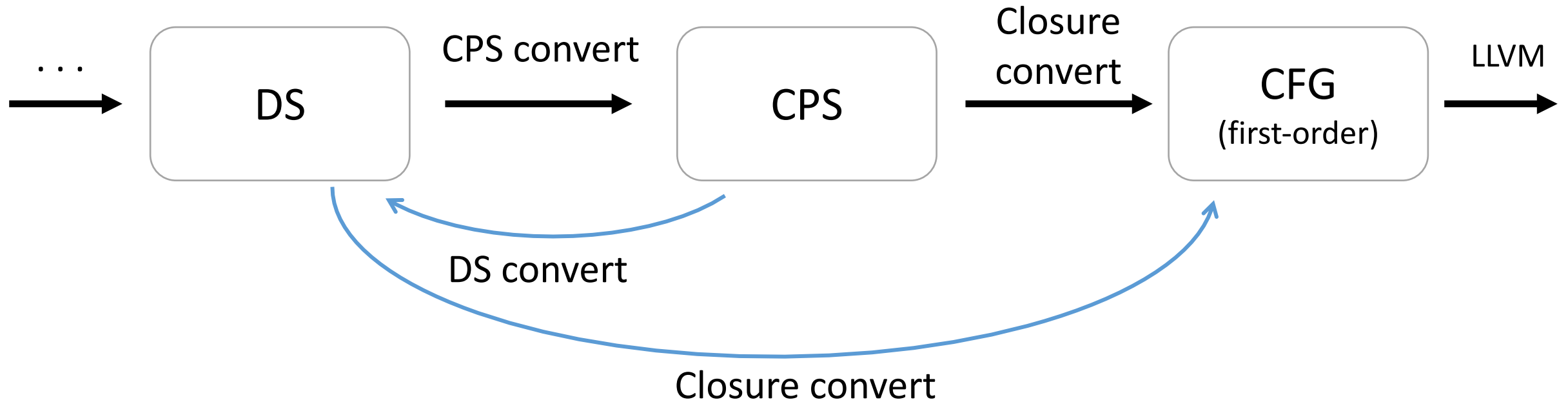
It starts with a good intermediate representation:

- Continuations and functions are different.
- Continuation parameters added by CPS are distinguished.

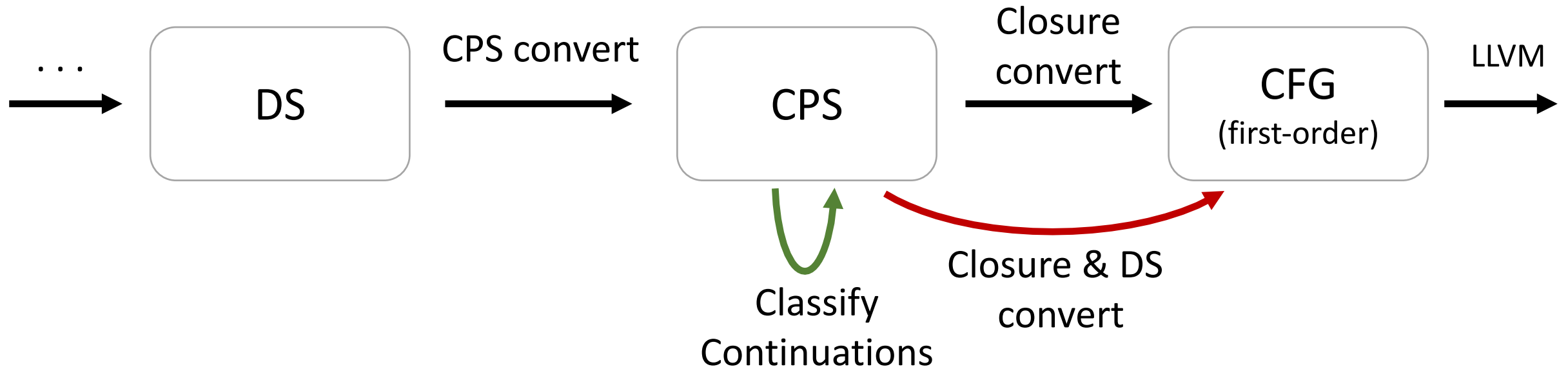
cont k () = _ **in** _ \leftrightarrow **throw** k ()

fun f (x, y / k) = _ \leftrightarrow f (1, 2 / k')

Noninvasive Compiler Upgrades



Noninvasive Compiler Upgrades



Classifying Continuations

Higher-order DS

```
fun g x = x  
fun f x y = if x > 10 then h((g x) + y) else h x
```

```
fun g (x / k) = throw k x ← Return throw
```

```
fun f (x, y / k) =
```

```
  cont doH z = h (z + y / k) in
```

```
  if x > 10
```

```
    then g (x / doH) ← Non-tail call
```

```
    else h (x / k) ← Tail call
```

Return continuations are only ever used or passed from the same function.

CPS

Converting to Direct Style

Higher-order CPS

```
fun g (x / k) = throw k x
fun f (x, y / k) =
  cont doH z = h (z + y / k) in
  if x > 10
  then g (x / doH)
  else h (x / k)
```

First-order DS

```
fun g (_, x) = return x
fun f (ep, x, y) =
  block doH (ep, z, y) =
    tailcall h (z + y)
  if x > 10
  then z = call g x
    goto doH (ep, z, y)
  else tailcall h x
```

Taming CPS Optimizations

- Arity raising
- Expansive inlining
- ... maybe others?

```
fun foo t = let  
  val x = #1(t)  
  val t = #2(t)  
  ...
```

Unbox tuple



```
fun foo x y = let  
  ...
```

Taming CPS Optimizations

```
fun foo (_ / fooRet) =  
  fun bar ( / barRet) = throw barRet ()
```

```
fun g(_ / gRet) =  
  if ...  
  then bar( / gRet)  
  else throw gRet ()
```

```
cont joinK () =  
  ...  
  throw fooRet ()  
in  
  g (_ / joinK)
```

Taming CPS Optimizations

```
fun foo (_ / fooRet) =  
  fun bar ( / barRet) = throw barRet ()  
  
  fun g(_ / gRet) =  
    if ...  
    then bar( / gRet)  
    else throw gRet ()  
  
cont joinK () =  
  ...  
  throw fooRet ()  
in  
  g (_ / joinK)
```

CFA says `barRet = joinK`,
so we inline the `throw` to `barRet`.

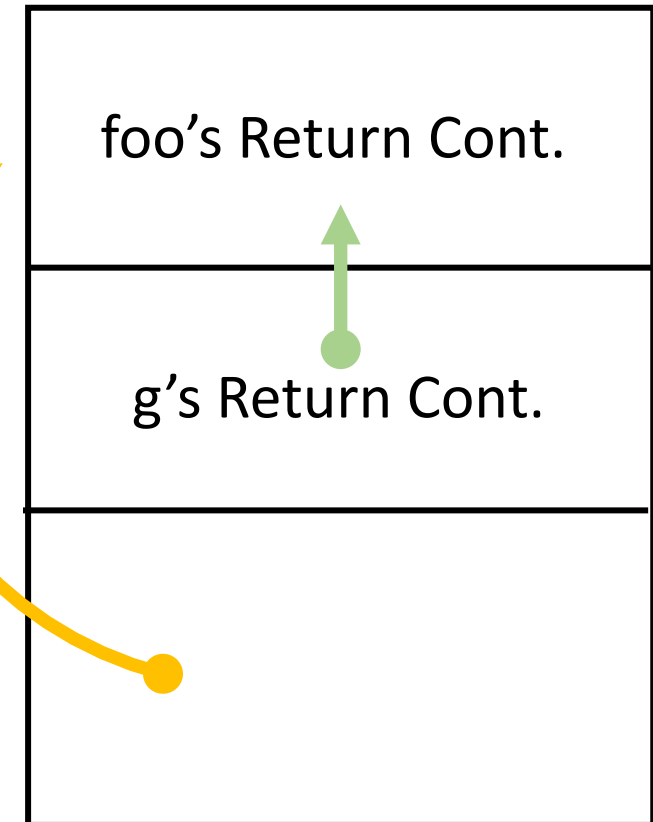
Taming CPS Optimizations

```
fun foo ( _ / fooRet ) =  
  fun bar ( / ) = ... throw fooRet ( )
```

```
fun g ( _ / gRet ) =  
  if ...  
  then bar ( / )  
  else throw gRet ( )
```

```
cont joinK ( ) =  
  ...  
  throw fooRet ( )  
in  
  g ( _ / joinK )
```

Stack (grows down)



Conclusion and Ongoing Work

- Direct style conversion can be done easily during closure conversion.
- Ongoing Work
 - Dedicated stack-based cont primitives (*newStack, etc.*)
 - Extending LLVM to support first-class stack-based conts.